

Everything Everywhere All at Once - How Improving Software Sustainability, Productivity, and Performance Naturally Go Hand in Hand

A story from the trenches of the US Exascale Computing Project

Prof. Dr. Hartwig Anzt

Technical University of Munich

05/23/2024



The US Exascale Computing Project



Advancing Scientific Discovery

The ECP aims to ensure availability of the exascale computing ecosystem necessary for developing clean energy systems, improving the resilience of our infrastructure, designing new materials that can perform in extreme environments, adapting to changes in the water cycle, developing smaller and more powerful accelerators for use in medicine and industry, and much more. Several projects focus on data-intensive problems to enable effective use of the data streams from powerful scientific facilities, complex environmental genomes, and cancer research (patient genetics, tumor genomes, molecular simulations, and clinical data).



Strengthening National Security

The ECP teams are also developing new applications for supporting the NNSA Stockpile Stewardship Program, which is responsible for maintaining the readiness and reliability of our nuclear weapons systems—without underground testing. Assessing the performance of weapons systems subject to hostile environments and potential threat scenarios exceeds the capabilities of current HPC systems and codes. NNSA application projects are focused on providing the sophisticated modeling and analysis tools needed to sustain the U.S. nuclear deterrence.



Improving Industrial Competitiveness

Exascale systems will be used to accelerate research that leads to innovative products and speeds commercialization, creating jobs and driving US competitiveness across industrial sectors, such as the emerging energy economy. To ensure alignment with US industry needs, the ECP is engaging senior technology decision makers from among the country's most prominent private sector companies.

The US Exascale Computing Project

Addressing a National Imperative

The Exascale Computing Project is an aggressive research, development, and deployment project focused on delivery of mission-critical applications, an integrated software stack, and exascale hardware technology advances.

Application Development



Software Technology

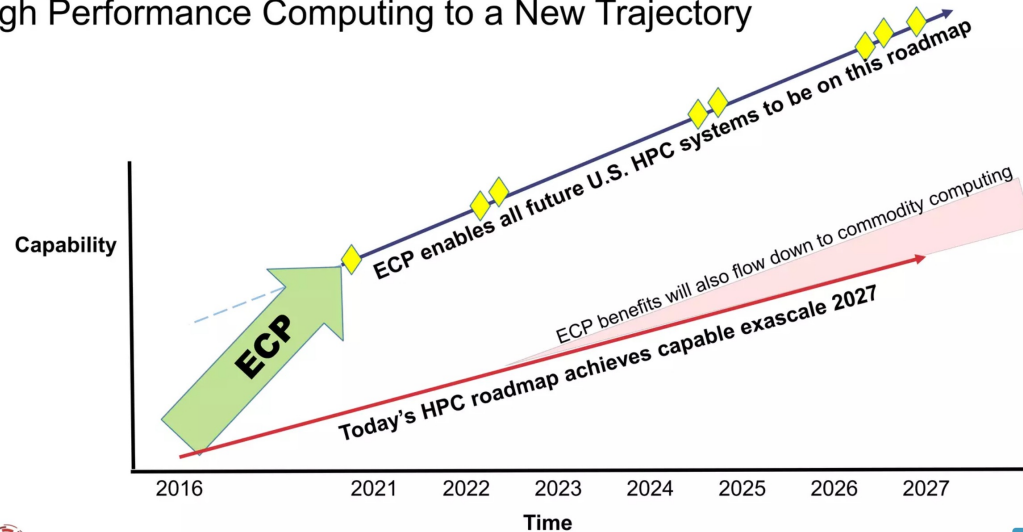


Hardware & Integration



© Paul Messina

Vision: Exascale Computing Project (ECP) Lifts all U.S. High Performance Computing to a New Trajectory



The US Exascale Computing Project



US\$4B – what is it spent on?

- 3 computers
 - \$600M each
 - \$400M to vendors for Design, Path, Fast - Forward 21 Applications



AMD Based
(Up & running)
20 MW



Intel Based
(Up & running)
40 MW



AMD APU Based
(planned)

The US Exascale Computing Project



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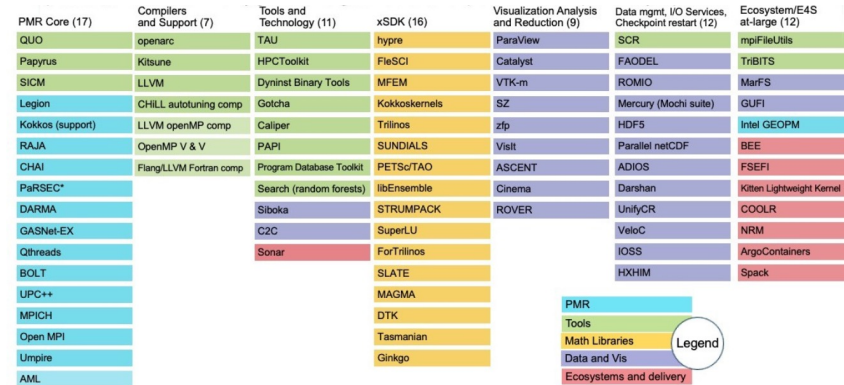


AMD APU Based
(panned)

Domain*	Base Challenge Problem
Wind Energy	2x2 5 MW turbine array in 3x3x1 km ³ domain
Nuclear Energy	Small Modular Reactor with complete in-vessel coolant loop
Fossil Energy	Burn fossil fuels cleanly with CLR's
Combustion	Reactivity controlled compression ignition
Accelerator Design	TeV-class 10 ²⁻³ times cheaper & smaller
Magnetic Fusion	Coupled gyrokinetics for ITER in H-mode
Nuclear Physics: QCD	Use correct light quark masses for first principles light nuclei properties
Chemistry: GAMESS	Heterogeneous catalysis: MSN reactions
Chemistry: NWChemEx	Catalytic conversion of biomass
Extreme Materials	Microstructure evolution in nuclear <u>matls</u>
Additive Manufacturing	Born-qualified 3D printed metal alloys

Domain*	Challenge Problem
Quantum Materials	Predict & control <u>matls</u> @ quantum level
Astrophysics	Supernovae explosions, neutron star mergers
Cosmology	Extract "dark sector" physics from upcoming cosmological surveys
Earthquakes	Regional hazard and risk assessment
Geoscience	Well-scale fracture propagation in wellbore cement due to attack of CO ₂ -saturated fluid
Earth System	Assess regional impacts of climate change on the water cycle @ 5 SYPD
Power Grid	Large-scale planning under uncertainty; underfrequency response
Cancer Research	Scalable machine learning for predictive preclinical models and targeted therapy
Metagenomics	Discover and characterize microbial communities through genomic and proteomic analysis
FEL Light Source	Protein and molecular structure determination using streaming light source data

Sustainable software development



A few words about myself

- Born and raised in Karlsruhe
- PhD in Numerical Mathematics from KIT
- Focus on computational linear algebra and high performance computing (HPC)
- Linear solvers, preconditioners, ...
- During my PostDoc at the University of Tennessee, I developed MAGMA sparse

MAGMA SPARSE

MAGMA-sparse as a “child” of MAGMA explores the development of sparse linear algebra functionality for NVIDIA GPUs.



Limitations:


- *C code with hand-written build system*
- *Sparse unit testing*
- *Focus on NVIDIA GPUs*
- *Design-specific limitations (flexibility/extensibility)*

Designing an ECP math library



Building Trusted Scientific Software

SHARE in f t p




PUBLISHED JUN 28, 2018 AUTHOR MIKE HER

I have worked in the scientific software field for more than 20 years. I often use the phrase "Verification is doing things right, and validation is doing things to make sure you are doing things right."

Pairing internal and external concerns
Verification focuses on internal concerns of a good software developer.

Software Verification

SHARE in f t p



PUBLISHED AUG 15, 2018 AUTHOR AMERU I

In the realm of software, verification is often erroneously used to mean the proper subset of verification for gaining confidence in the holistic process by which the developers convince themselves that the software they are developing is correct. It was designed to do. In scientific software this could mean numerical stability, and efficacy of the method in the expected results. Note that verification is limited to a model specification, not that the model itself matches the validation process.

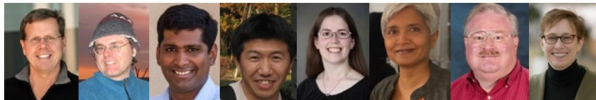
Think Locally, Act Globally: Outreach for Better Scientific Software

SHARE in f t p

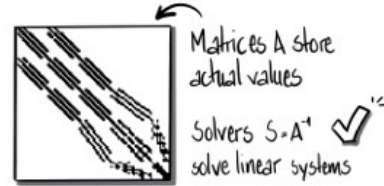
Helping code teams improve their software development, productivity, and sustainability is no small challenge. In the IDEAS Productivity project, we have found that one of the keys to aiding the Exascale Computing Project (ECP) software development teams involves extensive outreach to the broader community of computational scientists and engineers (CSE) in high-performance computing (HPC).

PUBLISHED JUL 17, 2018 AUTHOR DAVID BERNHOLDT TOPICS BETTER SKILLS PERSONAL PRODUCTIVITY AND SUSTAINABILITY

An ambitious goal
The ECP needs to deliver a software environment and applications ready to run on exascale computers, which are scheduled to be deployed starting in 2021. Achieving this goal entails a major, large-scale software development effort. Recognizing the challenges development teams will face, the ECP is supporting the IDEAS Productivity project to help scientific researchers improve their development practices.

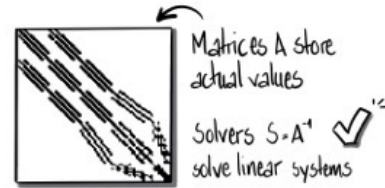
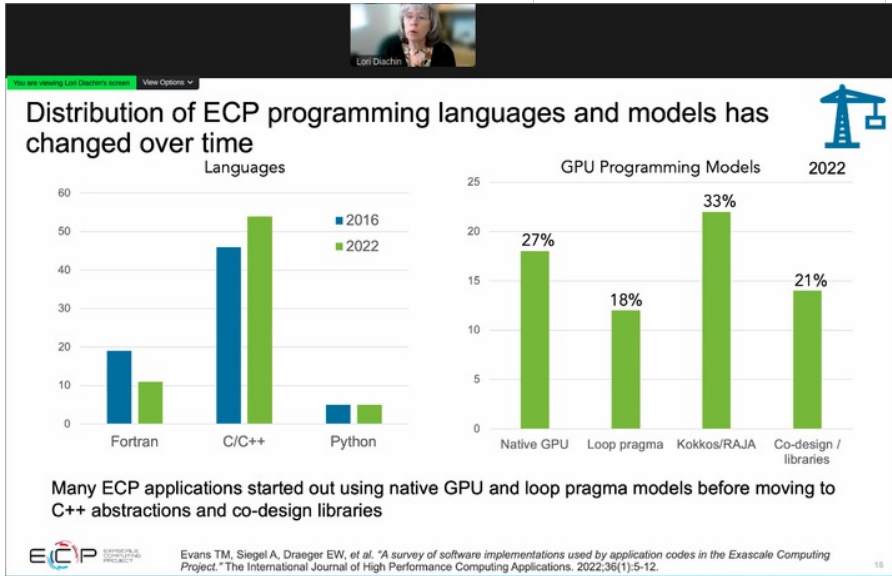


Ginkgo - A sparse linear algebra library for HPC



Designing software for performance, portability, & sustainability

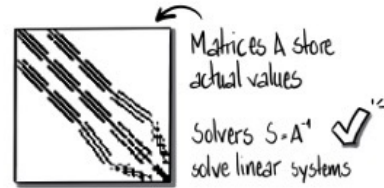
written in C++ → **Ginkgo** - A sparse linear algebra library for HPC



Lori Daichin, 05/22/2024

Designing software for performance, portability, & sustainability

written in C++ → **Ginkgo** - A sparse linear algebra library for HPC



Contains architecture-agnostic algorithm implementation

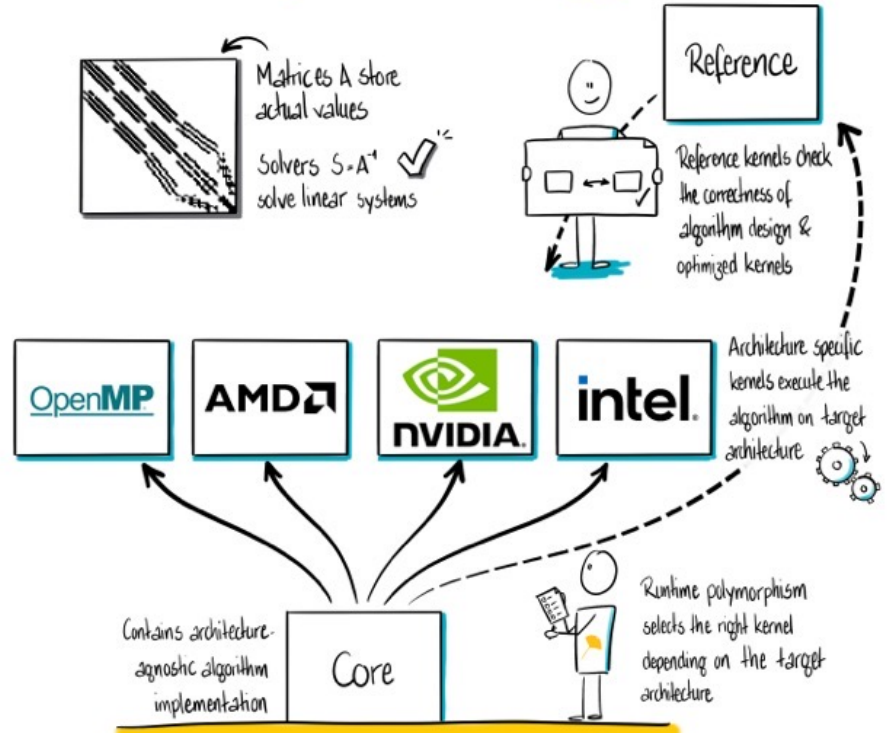
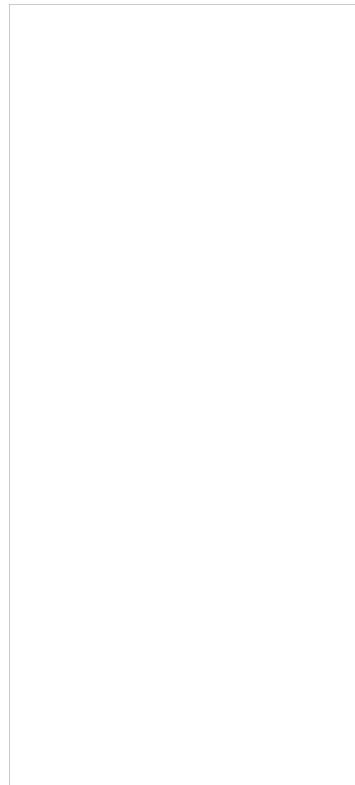
Core



Runtime polymorphism selects the right kernel depending on the target architecture

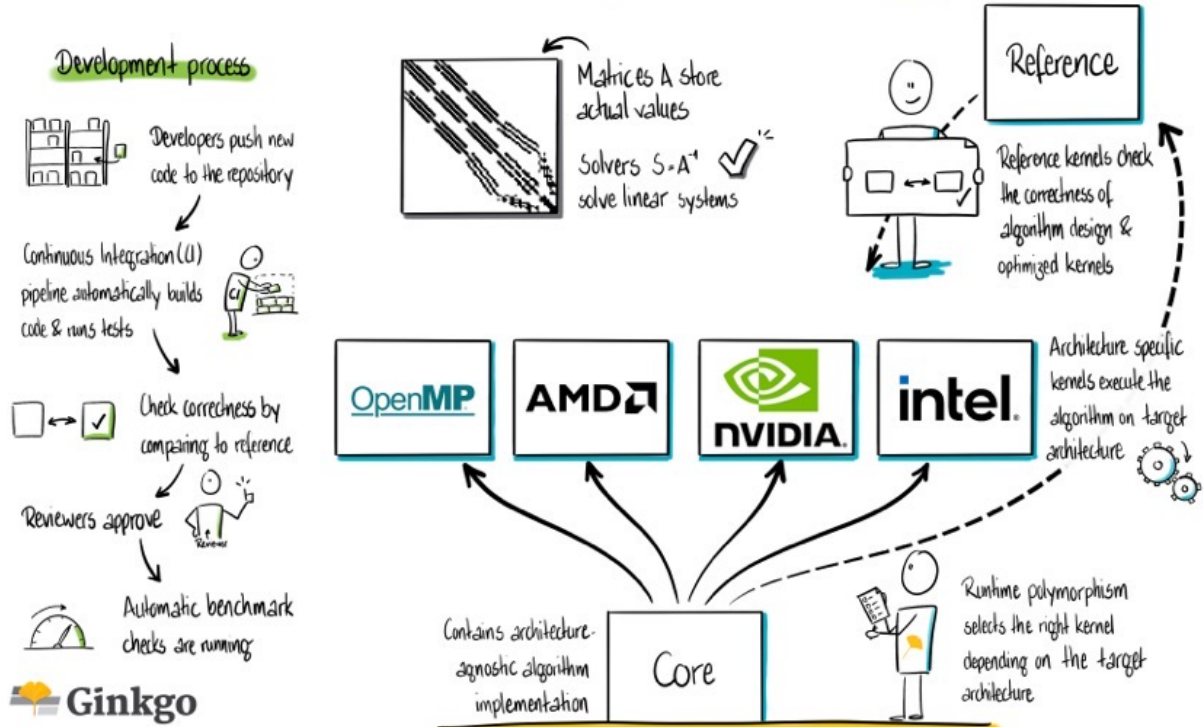
Designing software for performance, portability, & sustainability

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Designing software for performance, portability, & sustainability

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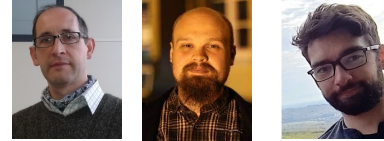


Ginkgo: A sparse linear algebra library for HPC

Linear Operator Interface

We express everything as Linear Operator.

- Internally, we leverage C++ class inheritance.
- Applications can apply any functionality as a linear operator.



Matrix-Vector Product

Preconditioner (for matrix A)

Solver (for system $Ax = b$)

$$x := A \cdot b$$

$$x := M^{-1} \cdot b$$

$$x := S \cdot b$$

$$M^{-1} \approx A^{-1}$$

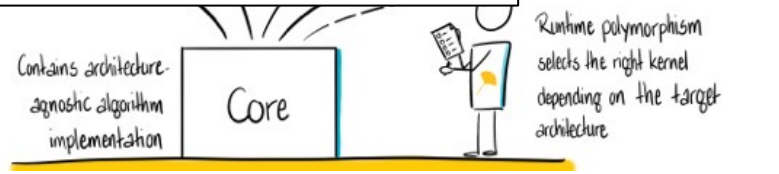
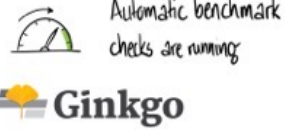
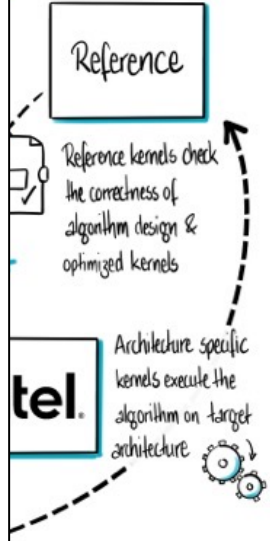
$$S \approx A^{-1}$$

$$M^{-1} = \Pi(A)$$

$$S = \Sigma(A)$$

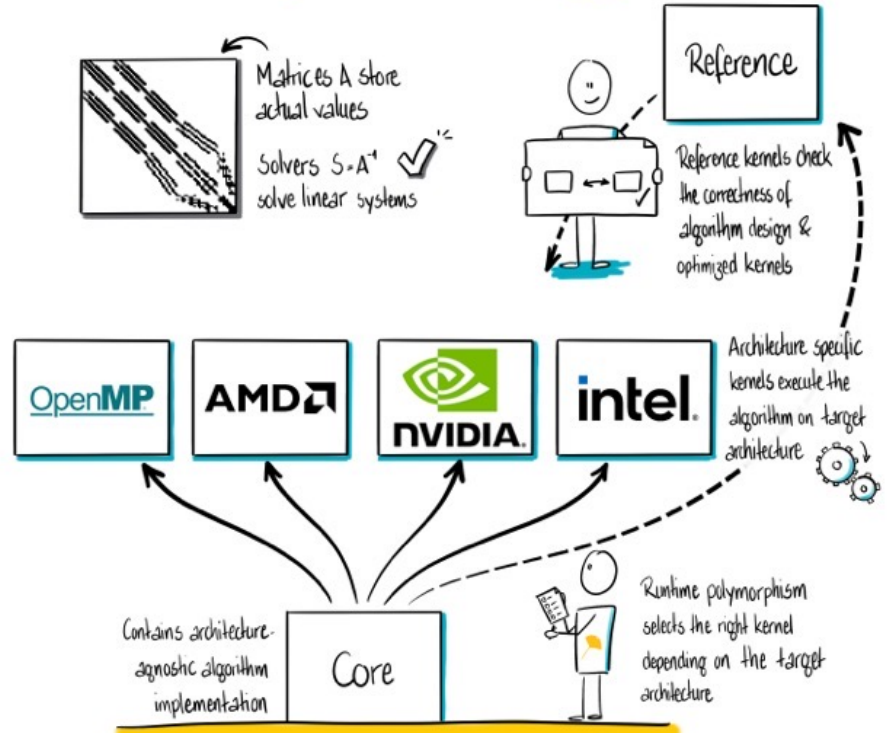
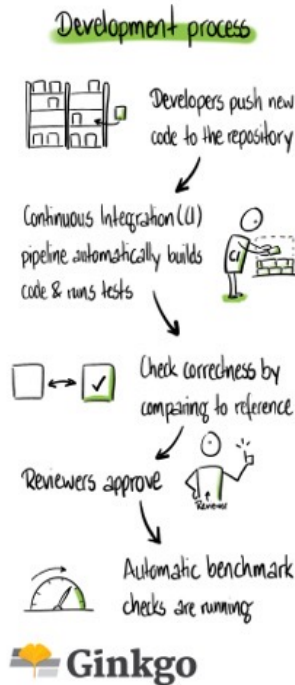
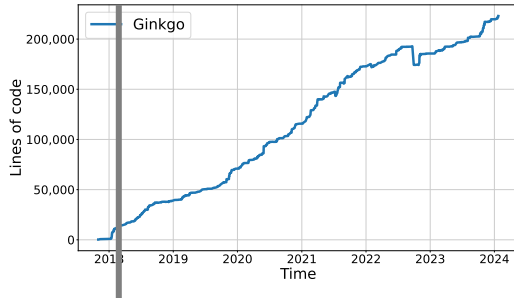
All of them can be expressed as

$$\text{Application of a linear operator* (LinOp)} \quad L : \mathbb{F}^m \rightarrow \mathbb{F}^m$$

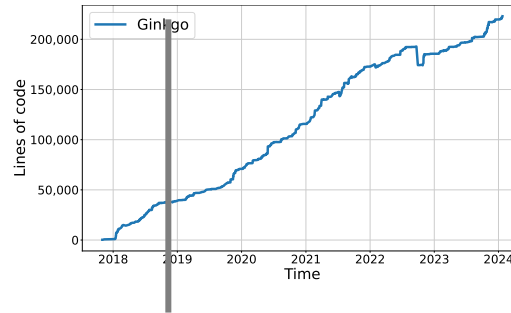
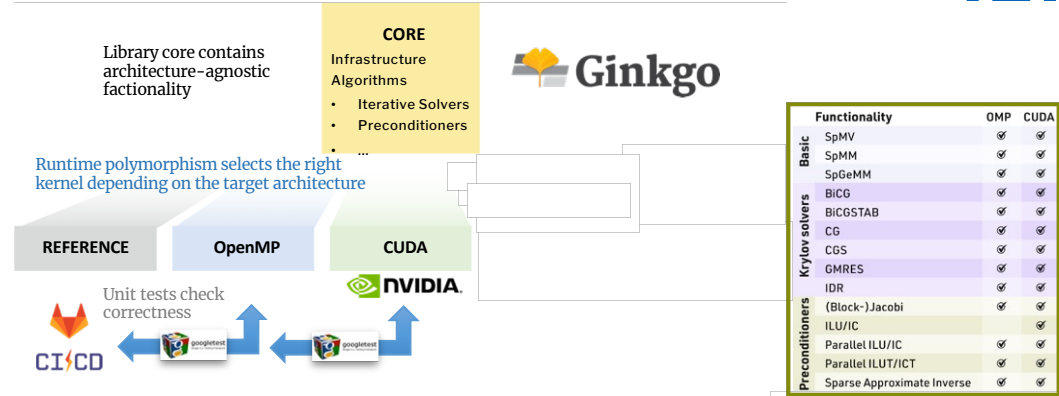


Designing software for performance, portability, & sustainability

written in C++  **Ginkgo** - A sparse linear algebra library for HPC



Starting with the CUDA backend



Extending to AMD GPUs



~2 months

better scientific software

Resources ▾ Blog Events About ▾

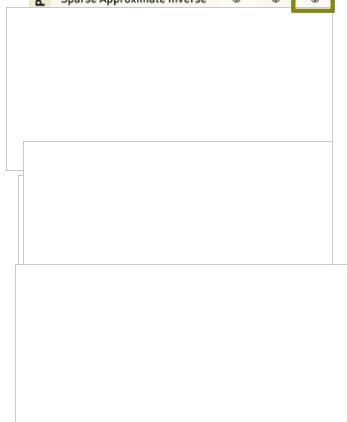
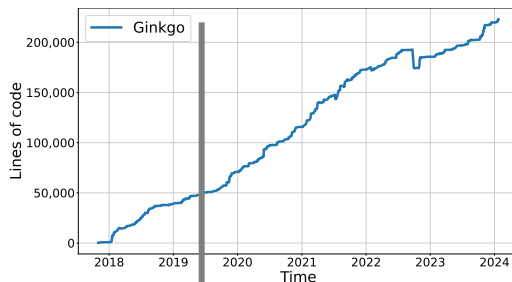
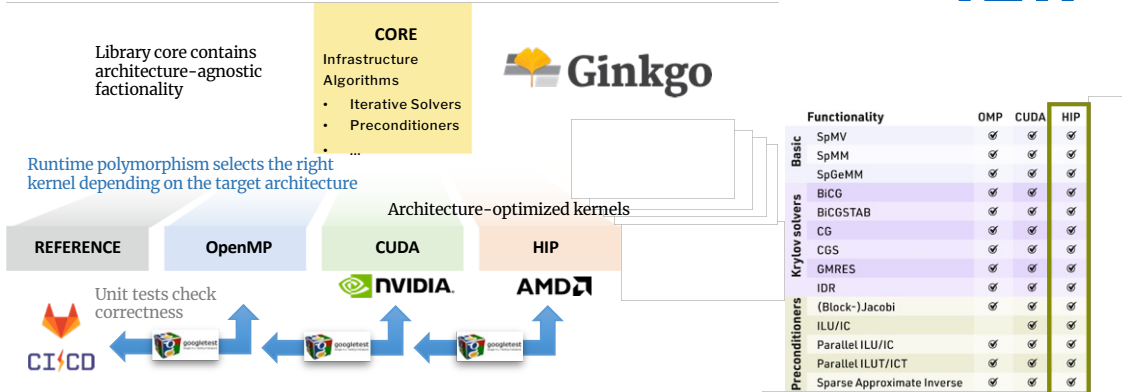
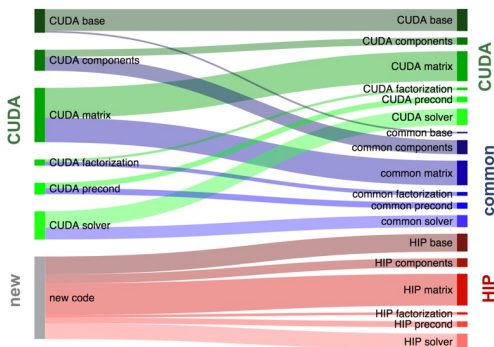
HOME > BLOG > Porting the Ginkgo Package to AMD's HIP...

Porting the Ginkgo Package to AMD's HIP Ecosystem

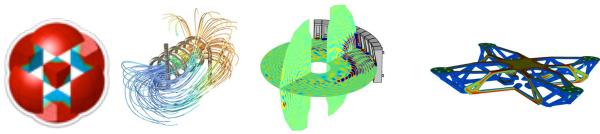
SHARE in f t

In response to the explosion-like diversification in hardware architectures, hardware portability and the ability to adopt new processor designs have become a central priority in realizing software sustainability. In this blog article, we discuss the experience of porting CUDA code to AMD's Heterogeneous-compute Interface for Portability (HIP).

PUBLISHED JUN 25, 2020 AUTHOR HARTWIG ANZT TOPICS BETTER RELIABILITY TESTING DESIGN



Input from the “first customer”



MFEM is a *free, lightweight, scalable* C++ library for finite element methods.

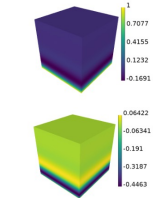
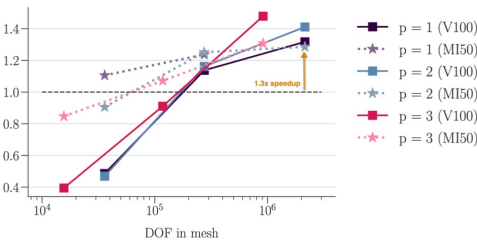
Speeding up MFEM’s “example 22” on GPUs

Example 22 of the MFEM finite element library solves harmonic oscillation problems, with a forced oscillation imposed at the boundary. In this test, we use variant 1:

$$-\nabla \cdot (a \nabla u) - \omega^2 b u + i \omega c u = 0$$

with $a = 1, b = 1, \omega = 10, c = 20$

Speedup for Ginkgo CB-GMRES vs MFEM



Real part of solution (top),
imaginary part of solution

Speedup of Ginkgo’s Compressed Basis-GMRES solver vs MFEM’s GMRES solver for three different orders of basis functions (p), using MFEM matrix-free operators and the Ginkgo-MFEM integration wrappers in MFEM. CUDA 10.1/V100 and ROCm 4.0/MI50.



Library core contains architecture-agnostic functionality

CORE
Infrastructure Algorithms

- Iterative Solvers
- Preconditioners

Runtime polymorphism selects the right kernel depending on the target architecture

Architecture-optimized kernels

REFERENCE

OpenMP

CUDA

HIP

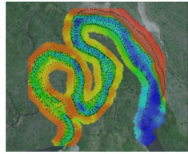
Unit tests check correctness

	OMP	CUDA	HIP
Basic			
SpMV	✓	✓	✓
SpMM	✓	✓	✓
SpGeMM	✓	✓	✓
BIG solvers			
BIGG	✓	✓	✓
BIGSTAB	✓	✓	✓
CG	✓	✓	✓
CGS	✓	✓	✓
GMRES	✓	✓	✓
IDR	✓	✓	✓
Preconditioners			
(Block-)Jacobi	✓	✓	✓
ILU/IC	✓	✓	✓
Parallel ILU/IC	✓	✓	✓
Parallel ILUT/ICT	✓	✓	✓
Sparse Approximate Inverse	✓	✓	✓

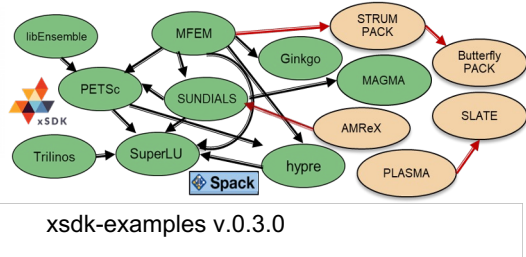
Utilities			
On-Device Matrix Assembly	✓	✓	✓
MC64/RCM reordering	✓		
Wrapping user data		✓	
Logging		✓	
PAPI counters		✓	

Part of the xSDK effort

xSDK: Extreme-scale Scientific Software Development Kit



Integrated surface-subsurface hydrology simulations of river meanders require the combined use of xSDK packages.



xsdk-examples v.0.3.0

The xSDK provides infrastructure for and interoperability of a **collection of related and complementary software elements**—developed by diverse, independent teams throughout the high-performance computing (HPC) community—that provide the building blocks, tools, models, processes, and related artifacts for rapid and efficient development of high-quality applications.

November 2022

- 26 math libraries
- 2 domain components
- 16 mandatory xSDK community policies
- Spack xSDK installer

xSDK community policies:

- 16 mandatory policies,
- 8 recommended policies,
- 4 Spack variant guidelines
- Available on Github <https://xsdk.info/policies/>

Library core contains architecture-agnostic functionality

Runtime polymorphism selects the right kernel depending on the target architecture

Architecture-optimized kernels

Unit tests check correctness

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CUDA

HIP

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Wrapping user data		✓	
Logging		✓	
PAPI counters		✓	

Extending to Intel GPUs



~12 months



Since 1987 - Covering the Fastest Computers in the World and the People Who Run Them

- Home
- Technologies
- Sectors
- COVID-19
- AI/ML/DL



March 23, 2021

yhmtsal/try_oneapi

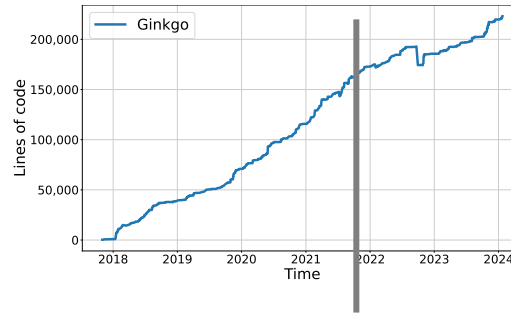
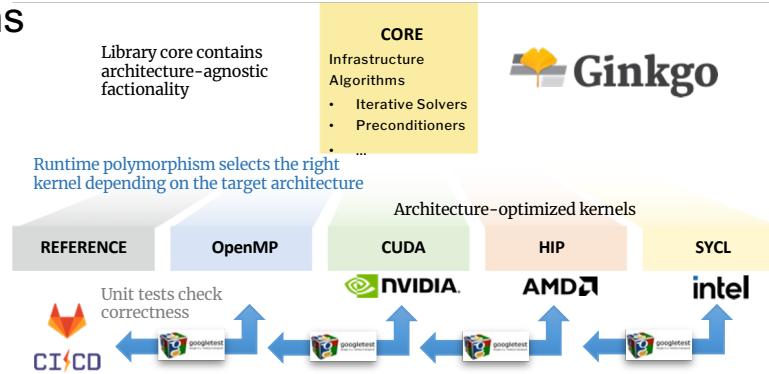
Code Issues Pull requests Actions Projects Security Insights

Go to file Add file Code About

yhmtsal format update 6/9/21/73 on Aug 7, 2022 70 commits

- arg_struct WIP 2 years ago
- atomic atomic and get_in_template 2 years ago
- check_unitin some checker last year
- classical_car gpu barrier issue in classical car spmv last year
- clinfo clinfo 2 years ago
- coop_cuda keep some history but I do not check them detail last year
- coop_draft keep some history but I do not check them detail last year

Releases
No releases published
Create a new release



Functionality	OMP	CUDA	HIP	DPC++
Basic				
SpMV	✓	✓	✓	✓
SpMM	✓	✓	✓	✓
SpGeMM	✓	✓	✓	✓
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Krylov solvers				
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Logging	✓	✓	✓	✓
PAPI counters	✓	✓	✓	✓

Extending to Intel GPUs



- Bi-Weekly technical meetings with Intel
- Long list of bug reports, feature requests, performance data discussions, documentation improvements ...

cuBLAS backend (and potentially other domains) fails with latest LLVM builds #223

mmetereel commented 22 days ago · edited · Contributor

Summary

As first observed in #19 many tests in cuBLAS backend is failing with latest LLVM builds.

Version

I have tried LLVM commit: 66361038b63caaa566c9648f1daf50b74222b83 and got the below tests failing (showing only a few of them)

- 1 - BLAS/RT/Non2TestSuite/Non2Tests.RealSinglePrecision/Column_Major_TITAN RTX (Failed)
- 3 - BLAS/RT/Non2TestSuite/Non2Tests.RealDoublePrecision/Column_Major_TITAN RTX (Failed)
- 17 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 17 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 19 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 23 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 27 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 36 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 36 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 81 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 81 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 85 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
- 85 - BLAS/RT/Non2TestSuite/Non2Tests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)

... but also docker image contributions and bug fixes!

tid % subgroup size >= 4 gives wrong division

(double) 1/a gives wrong result when the tid % subgroup size is 0. For example, when a = 1.07338829563753890 1/a should be 0.9316293125835232

If (local_id == assign_id) { a = double(1/a); } when assign_id < 4, Gen9 GPU still give the correct result when assign_id >= 4, Gen9 GPU gives wrong 0.9316293125835232. CPU has more worse result

It is connected to optimizations (not reproducible with O0). fp-speculation=off do not improve results. Ticket number: XDEPS-4031 ()

ginkgohub/oneapi.cuda11.6

DS/ARCH linux/amd64 COMPRESSED SIZE 6.63 GB LAST PUSHED 22 days ago by yhmatal

IMAGE LAYERS

- 1 ADD File ... in /
- 2 CMD ["bash"]
- 3 ENV NVIDIA_VISIBLE_DEVICES=""

fix cuda/hip backend location #219

ymhatal merged 2 commits into oneapi:stx-devel from yhmatal:fix_cuda_backend_location 20 days ago

Description

From intel/llvm#6407, it moves almost all headers from CLibcyl to sycl I followed #199 way make the header can use sycl* if they exist and allow the old intel lvm. I also update the CLibcyl.h which are not changed before.

All Submissions

- Do all unit tests pass locally? Attach a log. A: it is a compiling issue.
- Have you formatted the code using clang-format?

Bug fixes

- Have you added relevant regression tests. A: it is a compiling issue.
- Have you included information on how to reproduce the issue (either in a GitHub issue or in this PR)?

Reproduce: compile the latest intel lvm and this repo, it will not be able to compile due to missing headers.

ymhatal added 2 commits 2 months ago

- use the correct sycl path after intel/13#6407
- fix the missing sycl/sycl.hpp

mmetereel commented on Aug 1

@ymhatal Thanks for the PR. In the description 'from sycl/CL to CL, correct? My understanding is all header files moved

From [DPCPP AoT documentation](#), not clear:

- The options are also required at linking time? Unused in files without kernels?
- Any example of other projects integrating AoT in a CMake setup?

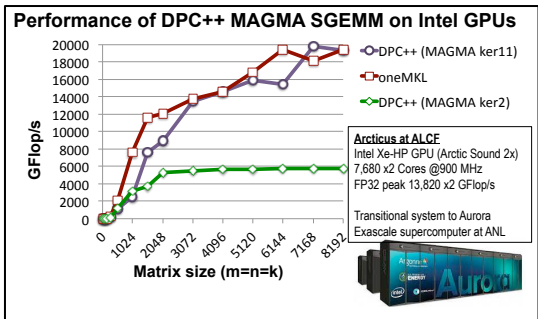
Intel Compiler (Fortran/C/C++/F03) - Intel Discrete GPU Accelerator - Joint Laboratory for System Evaluation (anl.gov)

hang_atomic_on_local

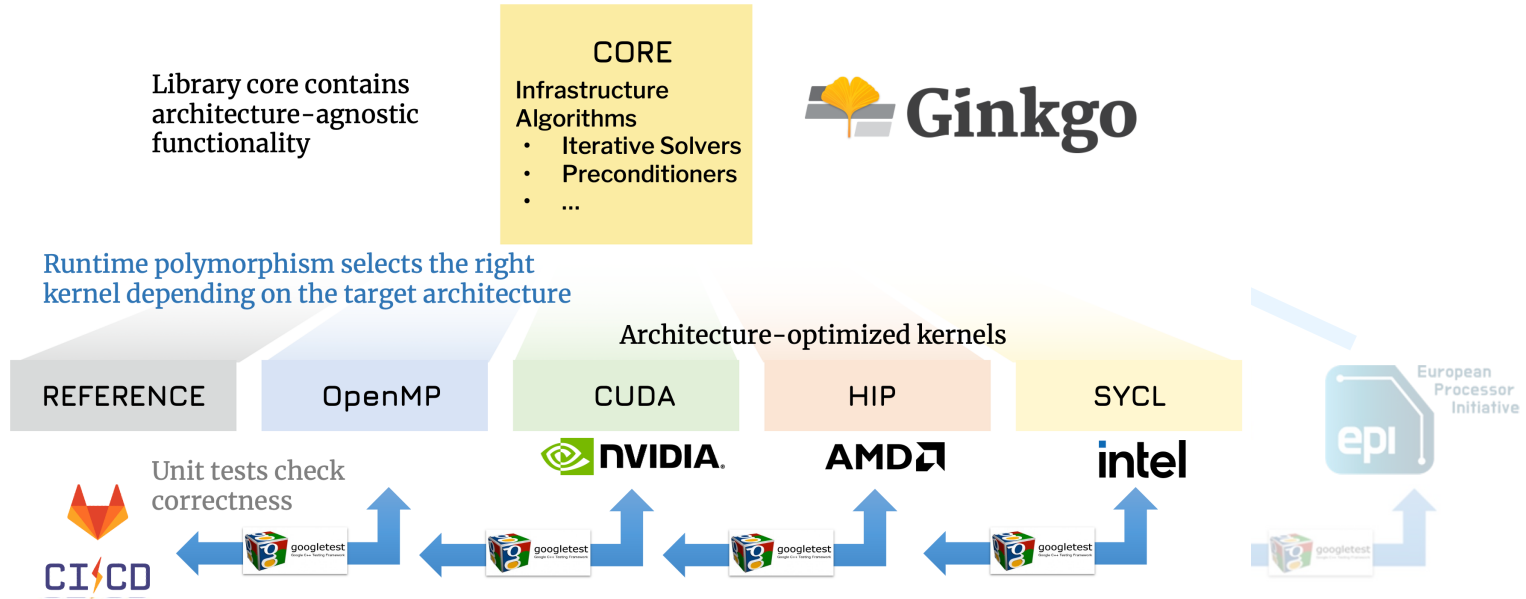
Ticket number: CMPLRLLVM-36572 (works in PVC, but still fails on ATS node) related to driver not compiler self

Jevcloud node issue

- sycl-ls/clinfo does not give any output
- no gpu on the nodes
- s001-n232, s001-n233, s011-n008
- nithub.com is not accessible on login



Portability as central design principle



This software design gives portability, performance, and sustainability.

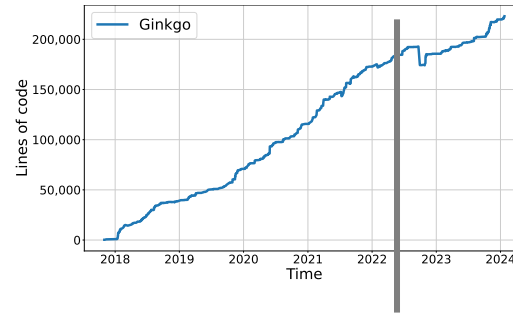
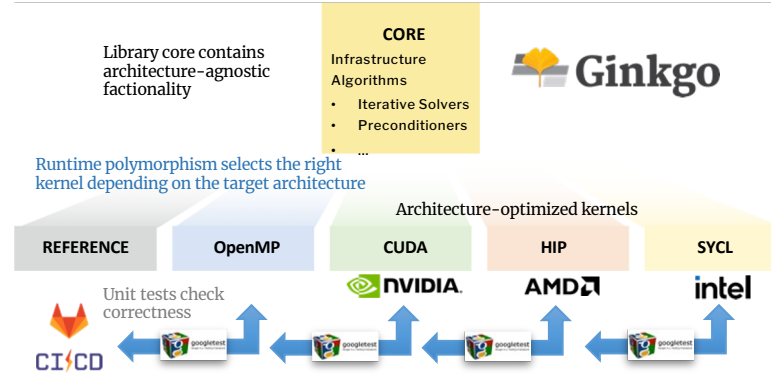
Focus efforts as lightweight tool in ECP to address challenges



Focus efforts



- Mixed precision
 - Address recent hardware trends (tensor cores, etc.)
- Batched Routines
 - Address application requirements



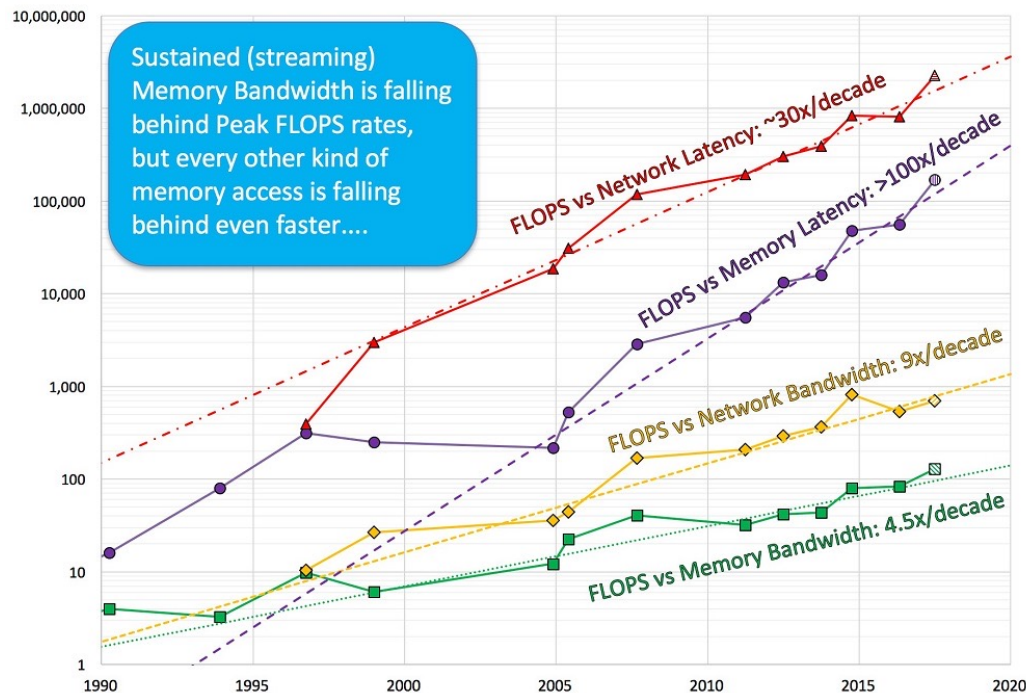
Functionality	OMP	CUDA	HIP	DPC++
Basic				
SpMV	✓	✓	✓	✓
SpMM	✓	✓	✓	✓
SpGeMM	✓	✓	✓	✓
Krylov solvers				
BICG	✓	✓	✓	✓
BICGSTAB	✓	✓	✓	✓
CG	✓	✓	✓	✓
CGS	✓	✓	✓	✓
GMRES	✓	✓	✓	✓
IDR	✓	✓	✓	✓
Preconditioners				
(Block-)Jacobi	✓	✓	✓	✓
ILU/IC	✓	✓	✓	✓
Parallel ILU/IC	✓	✓	✓	✓
Parallel ILUT/ICT	✓	✓	✓	✓
Sparse Approximate Inverse	✓	✓	✓	✓
Utilities				
On-Device Matrix Assembly	✓	✓	✓	✓
MC64/RCM reordering	✓			
Wrapping user data		✓		
Logging		✓		
PAPI counters		✓		

Mixed precision focus effort



Form Factor	H100 SXM
FP64	34 teraFLOPS
FP64 Tensor Core	67 teraFLOPS
FP32	67 teraFLOPS
TF32 Tensor Core	989 teraFLOPS ¹
BFLOAT16 Tensor Core	1,979 teraFLOPS ²
FP16 Tensor Core	1,979 teraFLOPS ²
FP8 Tensor Core	3,958 teraFLOPS ²
INT8 Tensor Core	3,958 TOPS ²
GPU memory	80GB
GPU memory bandwidth	3.35TB/s

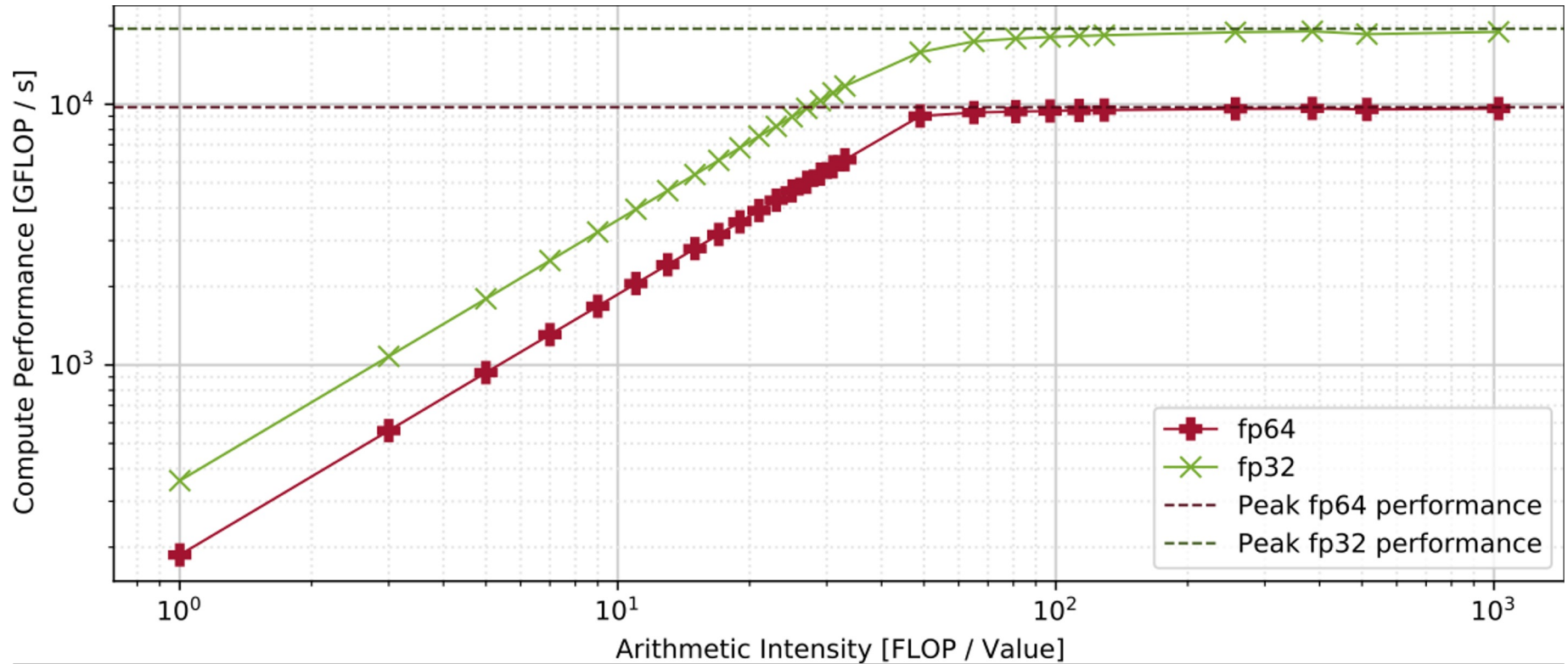
Balance: computation vs. communication



- (Dense) Matrix Performance
 - > Vector Operation Performance
- Low Precision Performance
 - > High Precision Performance

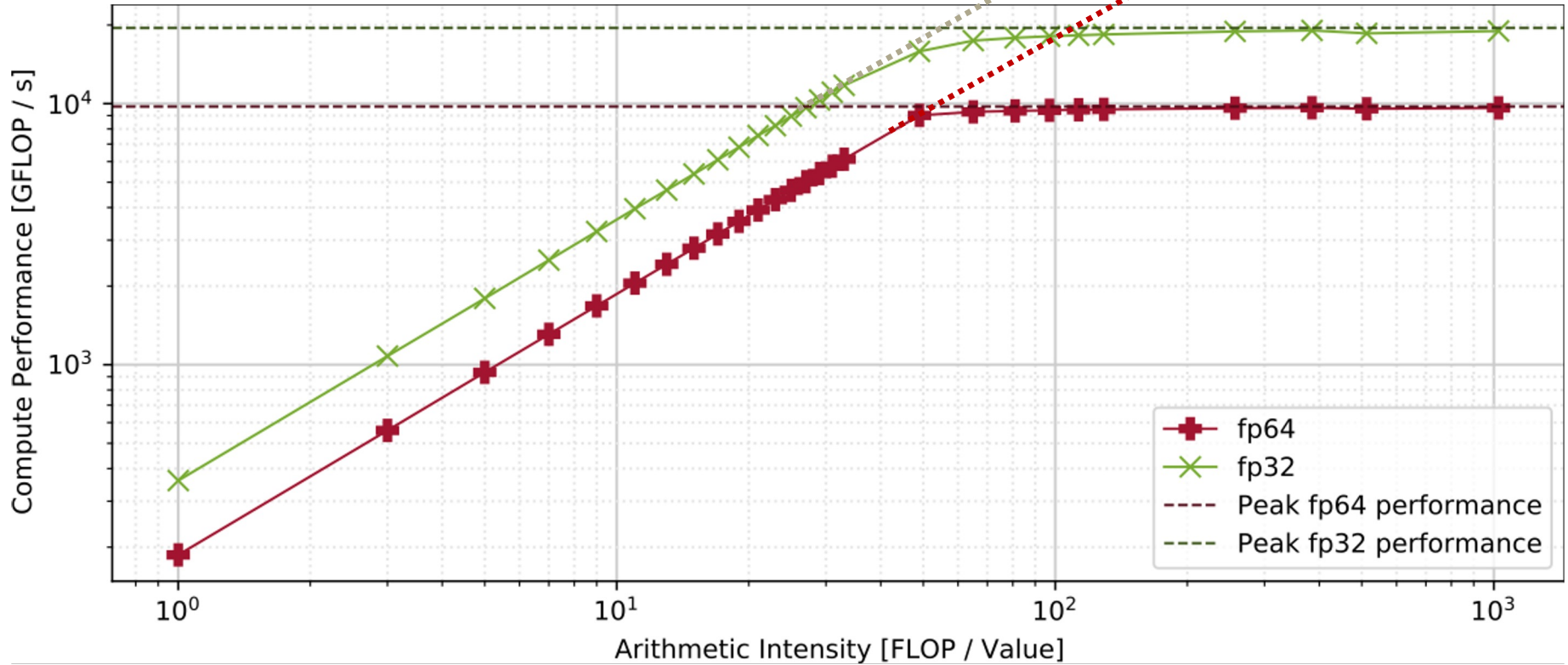
Trends in the relative performance of floating-point arithmetic and several classes of data access for select HPC servers over the past 25 years. Source: John McCalpin

NVIDIA A100



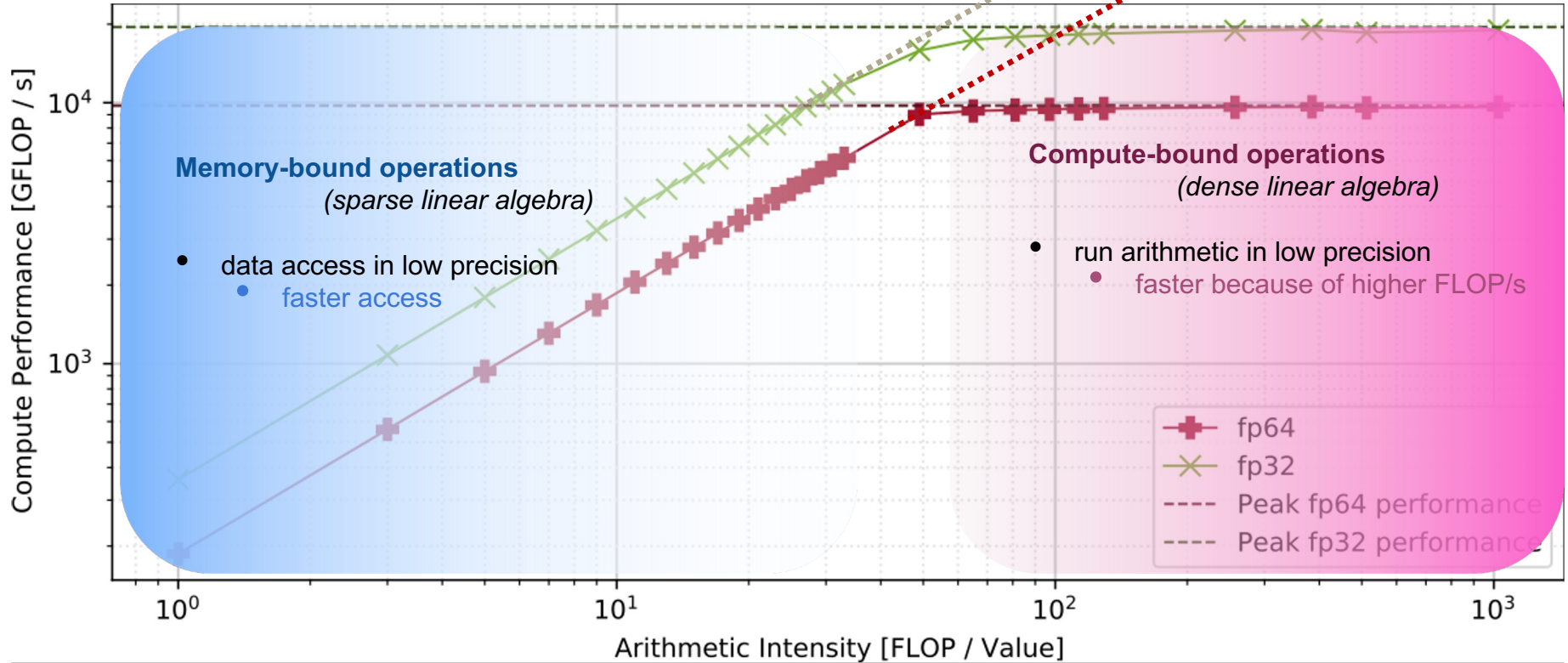
Matrix fp32

Matrix fp64



Matrix fp32

Matrix fp64

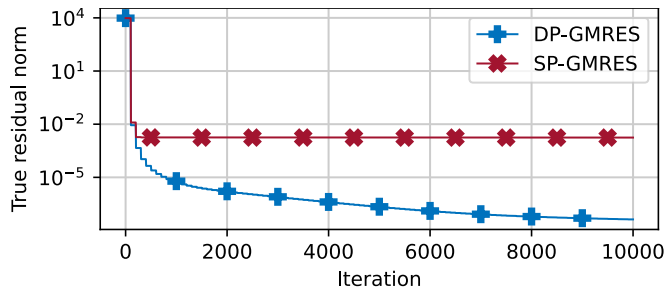


Linear System $Ax=b$ with $\text{cond}(A) \approx 10^7$
 (*apache2 from SuiteSparse*) **NVIDIA V100 GPU**

Double precision GMRES
 Initial residual norm **Relative residual $\sim 10^{-12}$**
 $\sqrt{r^t r}$: 9670.36
 Final residual norm
 $\sqrt{r^T r}$: $9.6639e-09$
 GMRES iteration count: 23271
 GMRES execution time: 43801 ms

Single precision GMRES
 Initial residual norm **Relative residual $\sim 10^{-7}$**
 $\sqrt{r^t r}$: 9670.36
 Final residual norm
 $\sqrt{r^T r}$: 0.00175464
 GMRES iteration count: 25000
 GMRES execution time: 27376 ms

$\sim 2x$ faster!

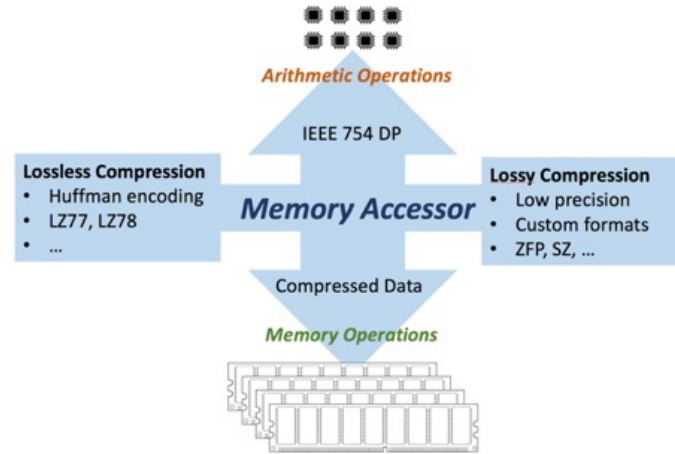


forward error \approx (unit round-off) * (linear system's condition number)

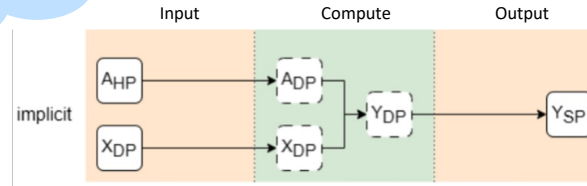
N. Higham: Accuracy and stability of numerical algorithms. SIAM, 2002.

Mixed precision focus effort

- Traditionally, we use a strong coupling between the precision formats used for **arithmetic operations** and **storing data**.
- *We should compute in fp64*
- *Data should be compressed for main memory access (low precision/compression)*
- *Compression / Conversion needs to happen on-the-fly*

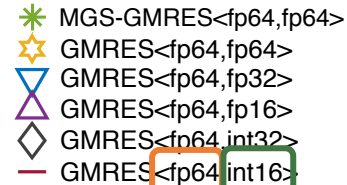


See Felix Liu's thesis

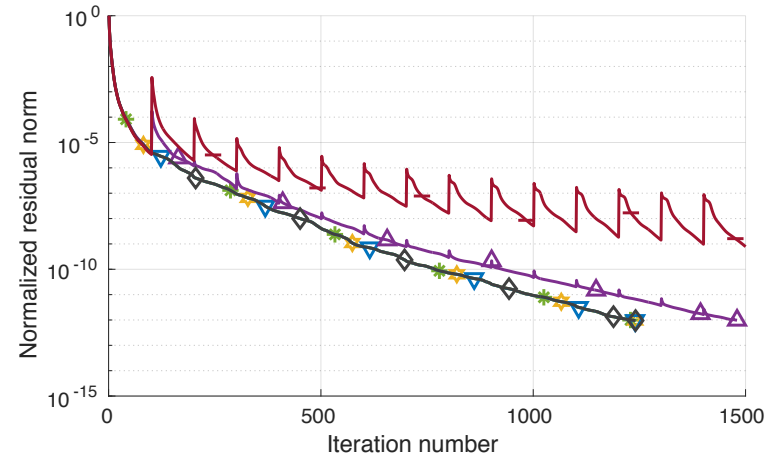


Compressed Basis (CB-) GMRES

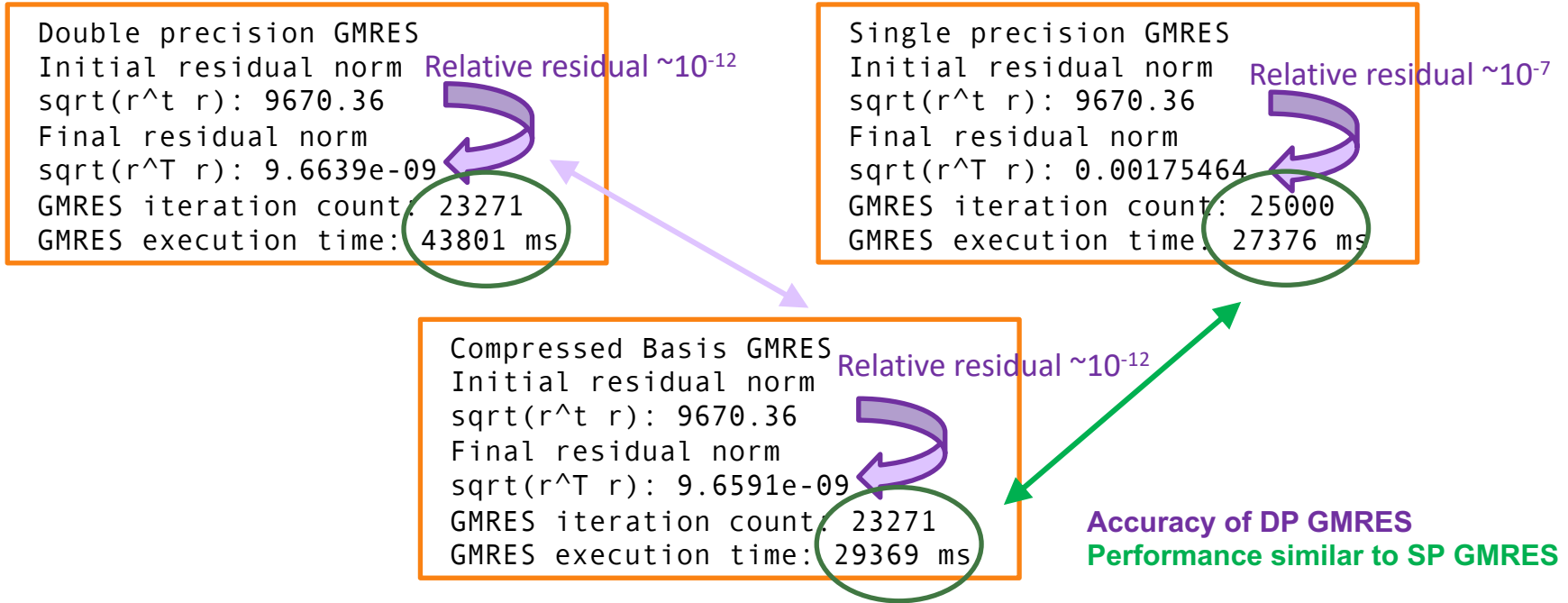
- Use double precision in all arithmetic operations;
- Store Krylov basis vectors in lower precision;
 - Search directions are no longer DP-orthogonal;
 - Hessenberg system maps solution to “perturbed” Krylov subspace;
 - Additional iterations may be needed;
 - As long as the loss-of-orthogonality is moderate, we should see moderate convergence degradation;

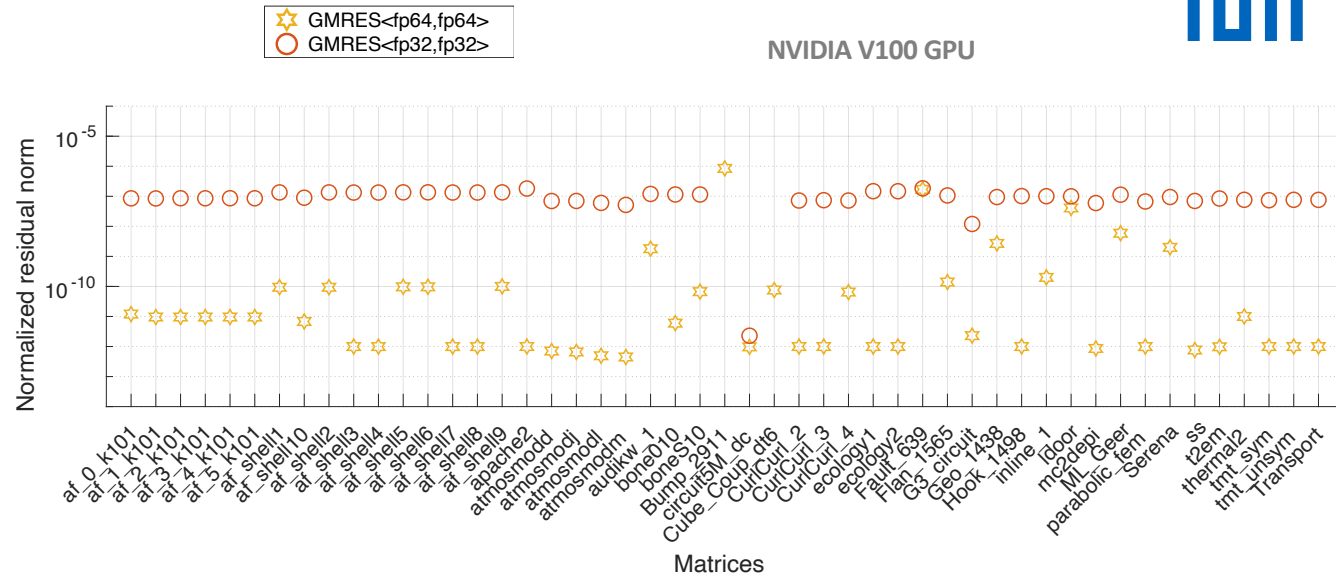


arithmetic precision memory precision

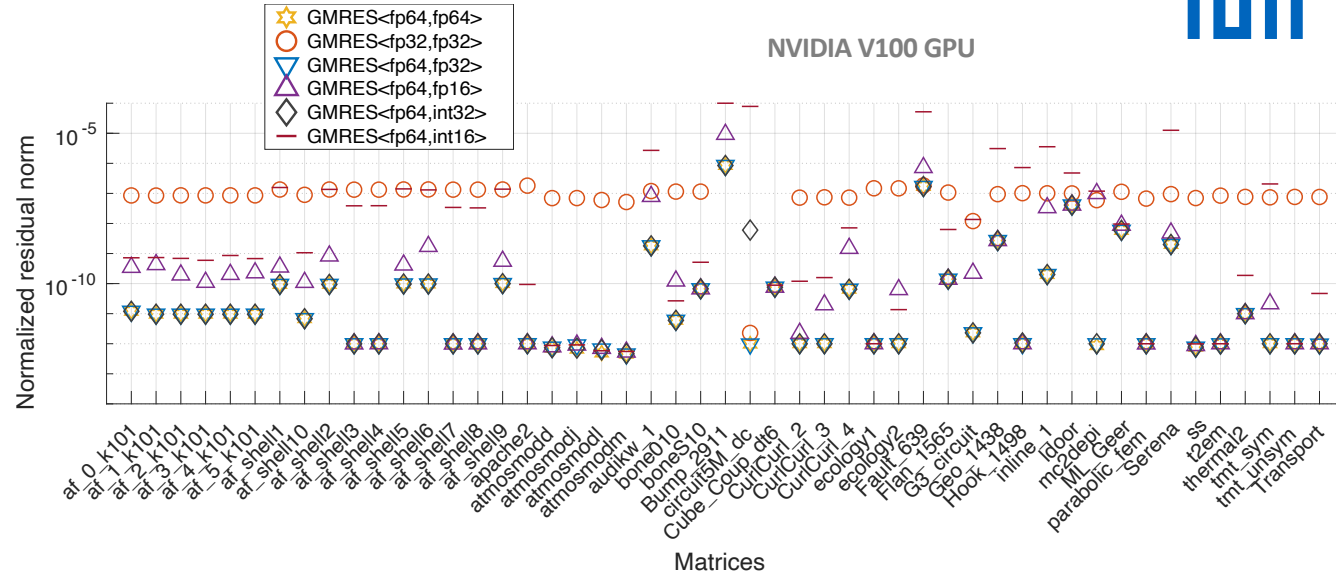


Linear System $Ax=b$ with $\text{cond}(A) \approx 10^7$
 (*apache2 from SuiteSparse*) **NVIDIA V100 GPU**





- CB-GMRES using 32-bit storage preserves DP accuracy (SP-GMRES does not)

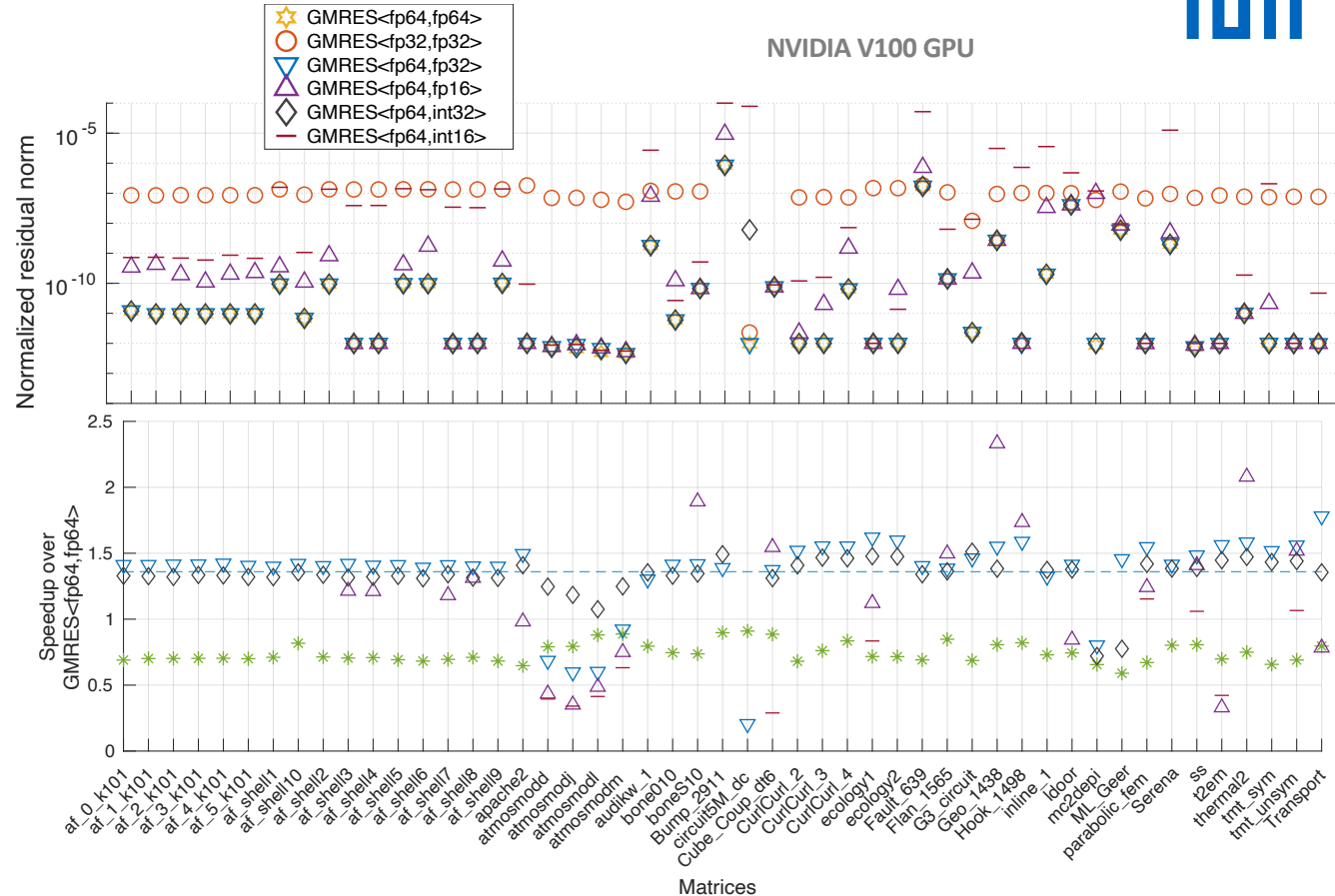


- CB-GMRES using 32-bit storage preserves DP accuracy (SP-GMRES does not)

- Speedups problem-dependent
- Speedup $\approx 1.4x$ (for restart 100)
- 16-bit storage mostly inefficient



Aliaga JI, Anzt H, Grützmacher T, Quintana-Ortí ES, Tomás AE. Compressed basis GMRES on high-performance graphics processing units. *The International Journal of High Performance Computing Applications*. 2022;0(0). doi:[10.1177/10943420221115140](https://doi.org/10.1177/10943420221115140)



Mixed precision AMG on GPUs

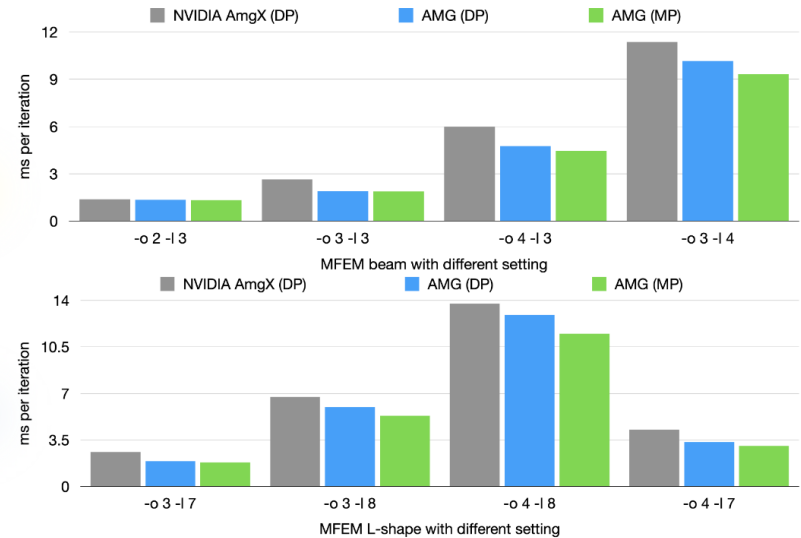
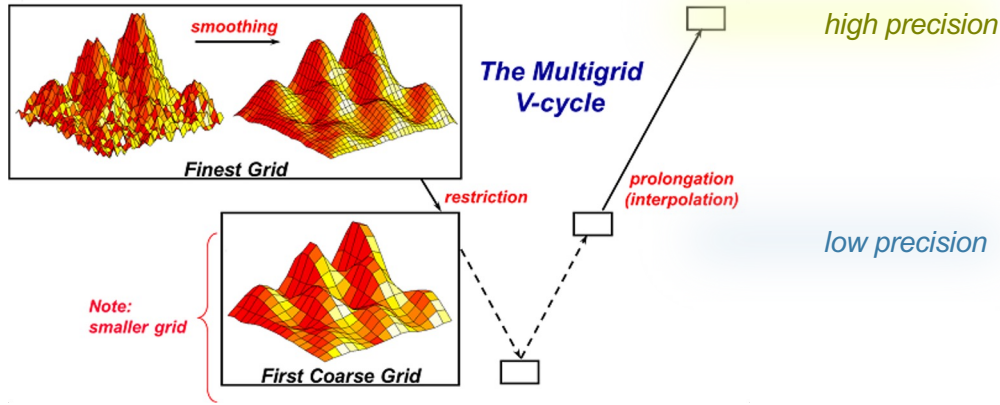


Mike Tsai

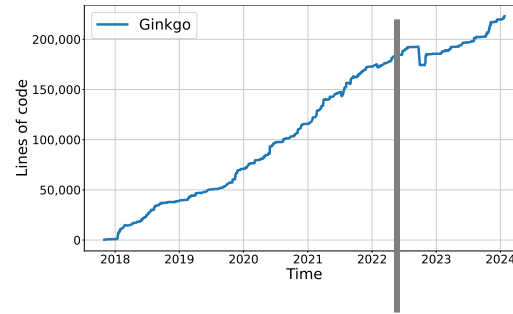
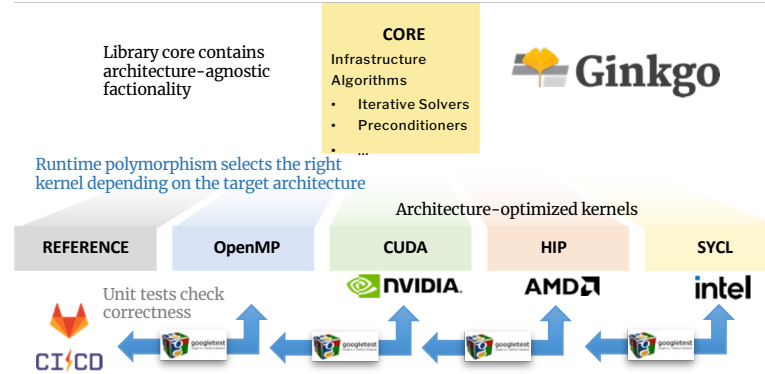
- **Preconditioning iterative solvers**

- Idea: Approximate inverse of system matrix to make the system “easier to solve”: $P^{-1} \approx A^{-1}$
and solve $Ax = b \Leftrightarrow P^{-1}Ax = P^{-1}b \Leftrightarrow \tilde{A}x = \tilde{b}$

- **Mixed Precision Multigrid Preconditioner**



Mixed precision focus effort



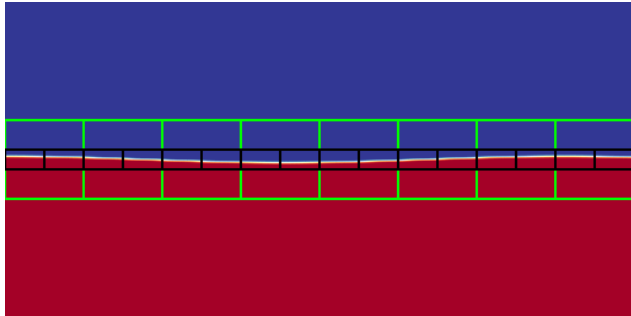
Functionality	OMP	CUDA	HIP	DPC++
Basic				
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SpMM	☑	☑	☑	☑
SpGeMM	☑	☑	☑	☑
BICG	☑	☑	☑	☑
BICGSTAB	☑	☑	☑	☑
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CGS	☑	☑	☑	☑
GMRES	☑	☑	☑	☑
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Krylov solvers				
(Block-)Jacobi	☑	☑	☑	☑
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Parallel ILUT/ICT	☑	☑	☑	☑
Sparse Approximate Inverse	☑	☑	☑	☑
AMG				
AMG preconditioner	☑	☑	☑	☑
AMG solver	☑	☑	☑	☑
Parallel Graph Match	☑	☑	☑	☑
Utilities				
On-Device Matrix Assembly	☑	☑	☑	☑
MC64/RCM reordering	☑			
Wrapping user data		☑		
Logging		☑		
PAPI counters		☑		

Batched focus effort – Combustion Simulations

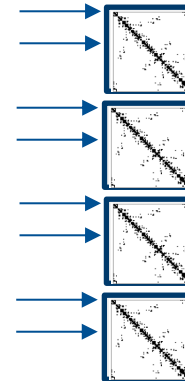
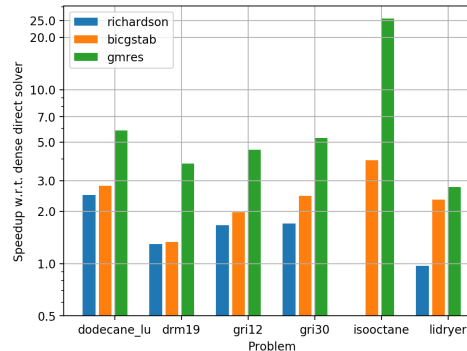
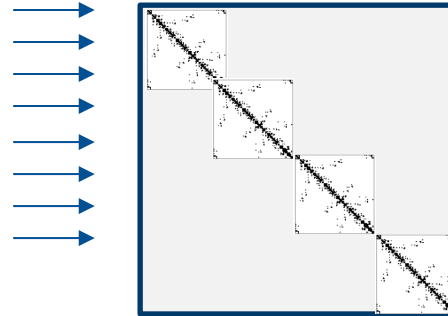
Batched iterative solvers for SUNDIALS / PeleLM

PeleLM is a parallel, adaptive mesh refinement (AMR) code that solves the reacting Navier-Stokes equations in the low Mach number regime. The core libraries for managing the subcycling AMR grids and communication are found in the [AMReX source code](https://amrex-combustion.github.io/AMReX_source_code).

<https://amrex-combustion.github.io/PeleLM/overview.html>



Problem	Size	Non-zeros (A)	Non-zeros (L+U)
dodecane_lu	54	2,332 (80%)	2,754 (94%)
drm19	22	438 (90%)	442 (91%)
gri12	33	978 (90%)	1,018 (93%)
gri30	54	2,560 (88%)	2,860 (98%)
isooctane	144	6,135 (30%)	20,307 (98%)
lidryer	10	91 (91%)	91 (91%)



Batched Sparse Iterative Solvers for Computational Chemistry Simulations on GPUs

Publisher: IEEE

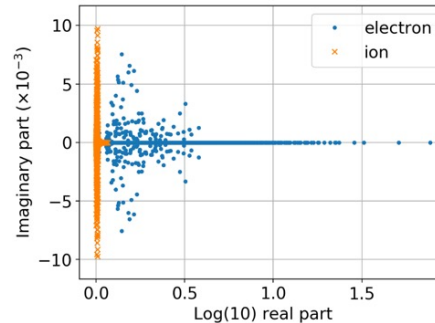
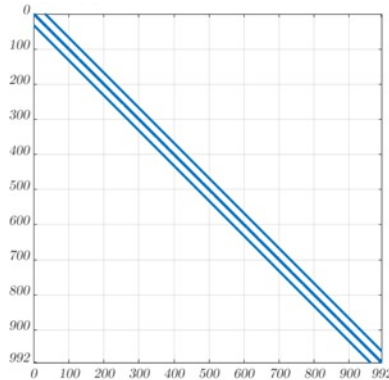
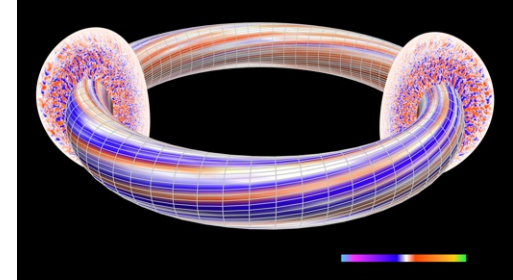
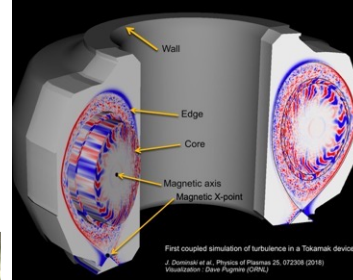
Cite This

PDF

Batched focus effort – Fusion Plasma Simulations

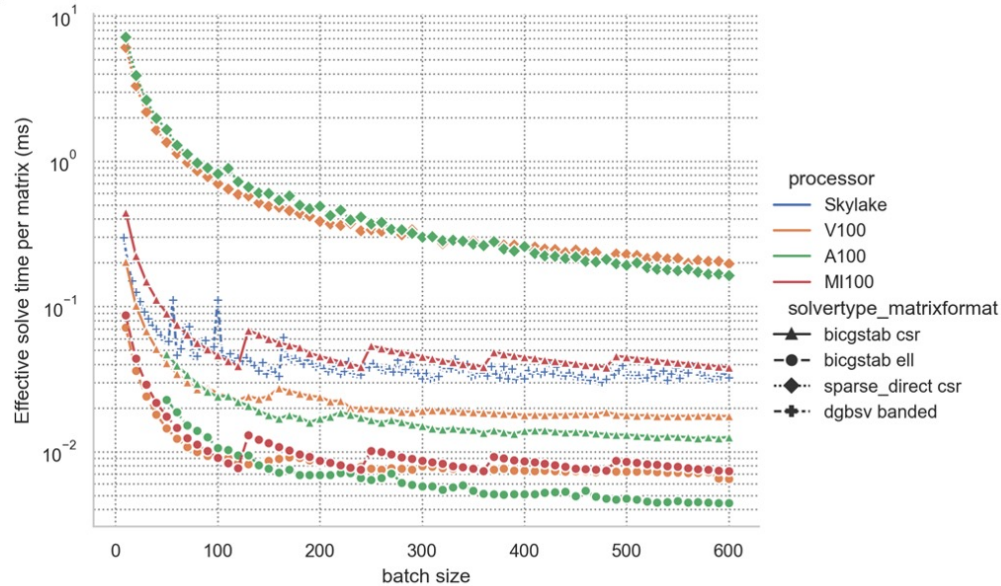
XGC is a gyrokinetic particle-in-cell code, which specializes in the simulation of the edge region of magnetically confined thermonuclear fusion plasma. The simulation domain can include the magnetic separatrix, magnetic axis and the biased material wall. XGC can run in total-delta-f, and conventional delta-f mode. The ion species are always gyrokinetic except for ETG simulation. Electrons can be adiabatic, massless fluid, driftkinetic, or gyrokinetic.

Source: https://xgc.pppl.gov/html/general_info.html

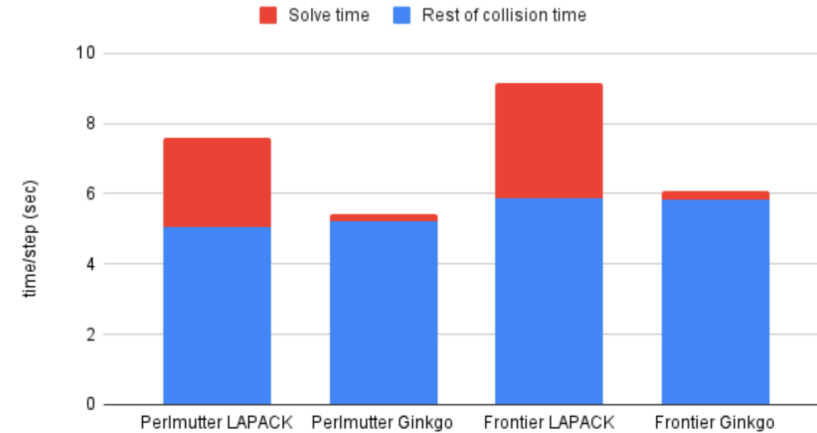


- Two species
- Ions easy to solve
- Electrons hard to solve
- Banded matrix structure
- Non-symmetric, need BiCGSTAB
- $n = \sim 1,000$
- $nz = \sim 9,000$

Batched focus effort – Fusion Plasma Simulations



XGC collision time reduction (64 nodes)



Aditya Kashi, Pratik Nayak, Dhruva Kulkarni, Aaron Scheinberg, Paul Lin, and Hartwig Anzt. **Batched sparse iterative solvers on gpu for the collision operator for fusion plasma simulations.** In *2022 IEEE International Parallel and Distributed Processing Symposium (IPDPS)*, pages 157–167. IEEE, 2022.

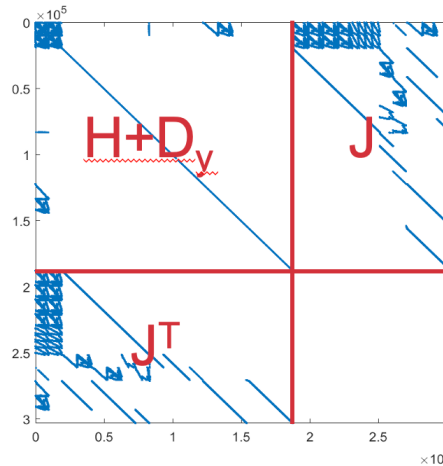
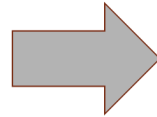


Underlying KKT Linear System Properties

See Felix Liu's thesis

- Security constrained optimal power flow analysis.
- The interior method strategy leads to symmetric indefinite linear systems

$$\underbrace{\begin{bmatrix} H + D_y & J \\ J^T & 0 \end{bmatrix}}_{K_k} \underbrace{\begin{bmatrix} \Delta y \\ \Delta \lambda \end{bmatrix}}_{\Delta x_k} = \underbrace{\begin{bmatrix} r_y \\ r_\lambda \end{bmatrix}}_{r_k},$$

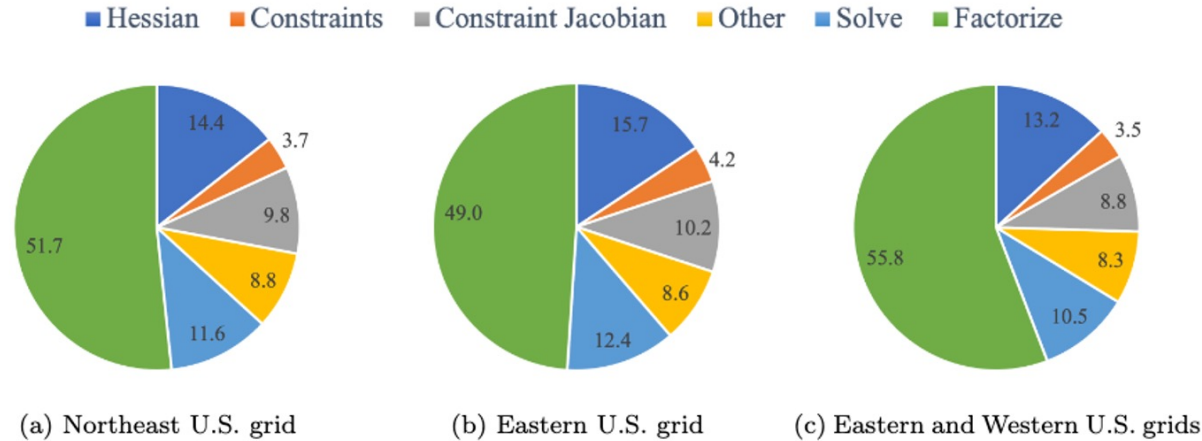


Typical sparsity pattern of optimal power flow matrices: No obvious structure that can be used by linear solver.

- The challenge: we need to solve a long sequences of such systems.

Sparse direct solvers for power grid simulations

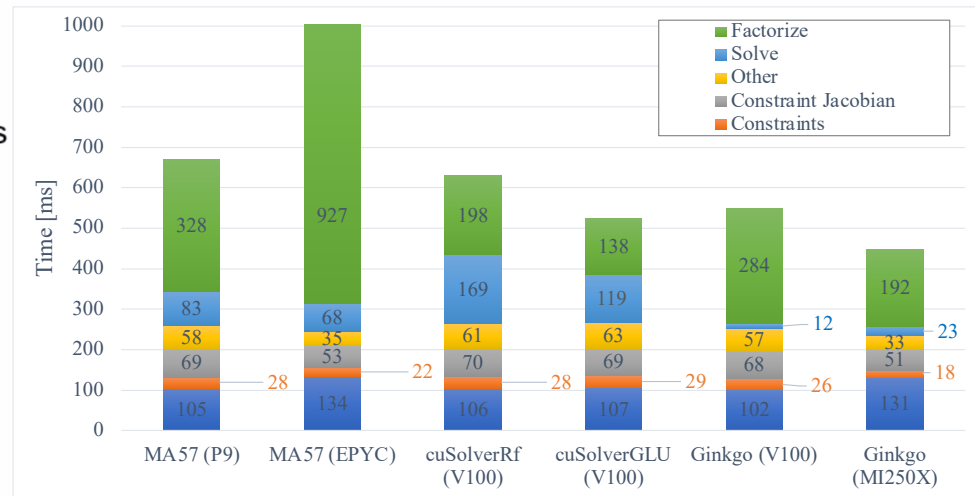
Grid	Buses	Generators	Lines	$N(K_k)$	$\text{nnz}(K_k)$
Northeastern US	25 K	4.8 K	32.3 K	108 K	1.19 M
Eastern US	70 K	10.4 K	88.2 K	296 K	3.20 M
Western and Eastern US	82 K	13.4 K	104.1 K	340 K	3.73 M



Sparse direct solvers for power grid simulations

Liner Solver Performance within Optimization Algorithm Average per iteration times (including first iteration on CPU)

- Each GPU solution outperforms all CPU baselines.
- Ginkgo performance improves on a better GPU.
- Iterative refinement configuration affects linear solver performance and optimization solver convergence.
- Ginkgo is the first GPU-resident sparse direct linear solver.



Now, after the completion of ECP

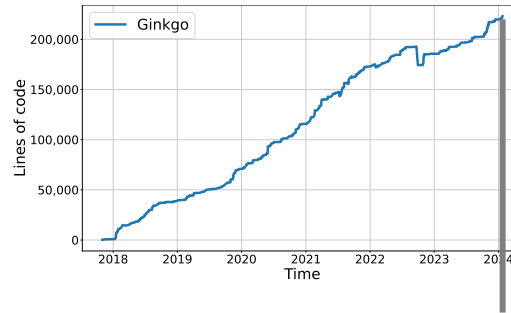
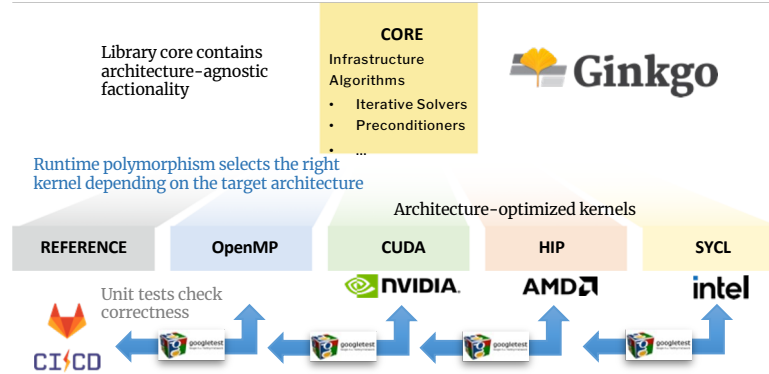


- Sustainable software design ready for the addition of new backends.
- EuroHPC Project MICROCARD uses Ginkgo

<https://www.microcard.eu>

- BMBF PDExa and ExaSIM projects use Ginkgo

- Companies are evaluating Ginkgo



FUNCTIONALITY		OMP	CUDA	HIP	DPC+
Basic	SpMV	✓	✓	✓	✓
	SpMM	✓	✓	✓	✓
	SpGeMM	✓	✓	✓	✓
	BICG	✓	✓	✓	✓
	BICGSTAB	✓	✓	✓	✓
Krylov solvers	CG	✓	✓	✓	✓
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	GCR	✓	✓	✓	✓
	GMRES	✓	✓	✓	✓
	FCG	✓	✓	✓	✓
	FGMRES	✓	✓	✓	✓
	IR	✓	✓	✓	✓
	IDR	✓	✓	✓	✓
	Block-Jacobi	✓	✓	✓	✓
	ILU/IC	✓	✓	✓	✓
Preconditioners	Parallel ILU/IC	✓	✓	✓	✓
	Parallel ILUT/ICT	✓	✓	✓	✓
	ISAI	✓	✓	✓	✓
	Batched BICGSTAB	✓	✓	✓	✓
	Batched CG	✓	✓	✓	✓
Batched	Batched GMRES	✓	✓	✓	✓
	Batched ILU	✓	✓	✓	✓
	Batched ISAI	✓	✓	✓	✓
	Batched Block-Jacobi	✓	✓	✓	✓
	AMG preconditioner	✓	✓	✓	✓
AMG	AMG solver	✓	✓	✓	✓
	Parallel Graph Match	✓	✓	✓	✓
	Symbolic Cholesky	✓	✓	✓	✓
	Numeric Cholesky	✓	✓	✓	✓
	Symbolic LU	✓	✓	✓	✓
Sparse direct	Numeric LU	✓	✓	✓	✓
	Sparse TRSV	✓	✓	✓	✓
	On-Device Matrix Assembly	✓	✓	✓	✓
	MC64/RCM reordering	✓	✓	✓	✓
	Wrapping user data	✓	✓	✓	✓
Utilities	Logging	✓	✓	✓	✓
	PAPI counters	✓	✓	✓	✓

MPI Support Single-GPU Support

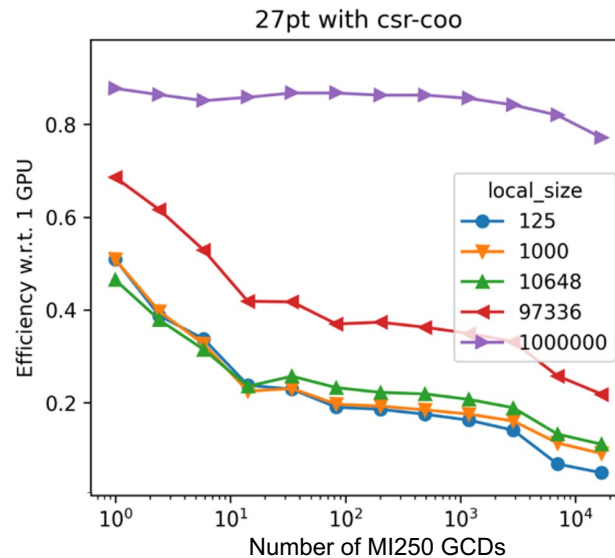
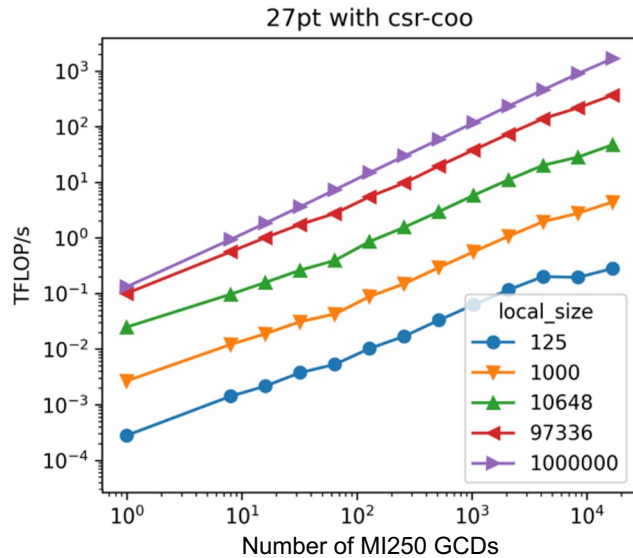


Scalability of Ginkgo on Frontier (#1 TOP500, AMD MI250)



Weak scaling: problem size increases with parallel resources

Weak scaling up to 8k AMD MI250 GPUs (16k GCDs)



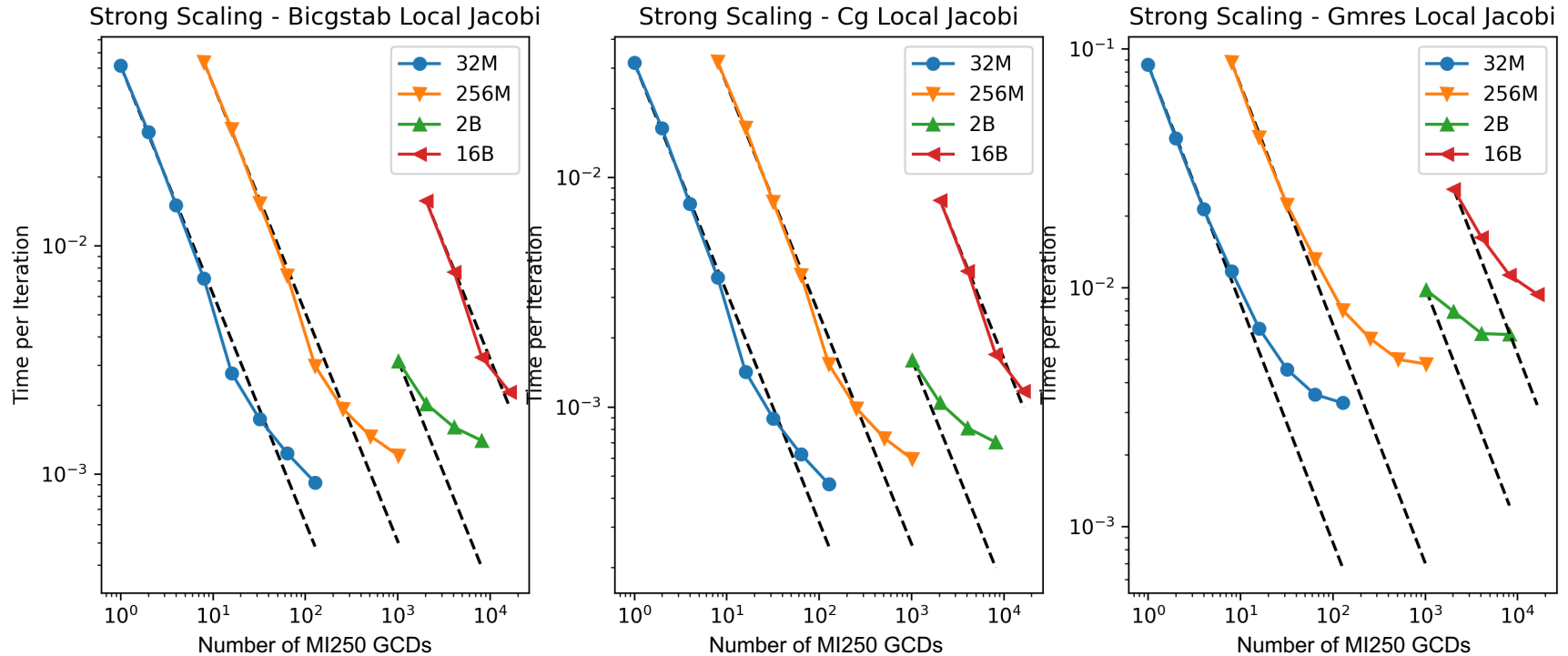
Significant Compute Waste!

Scalability of Ginkgo on Frontier (#1 TOP500, AMD MI250)



Strong scaling: problem size increases with parallel resources

Strong scaling up to 8k AMD MI250 GPUs (16k GCDs)



Lessons learnt over the years



- **ECP earmarking roughly half the budget to Software & App development is a game changer.**
 - **Central component for the success of ECP.**
 - This concept needs to – and does become - the blueprint for other nations, companies, and projects.
- **Workforce recruitment and workforce retention are the key to success in software development.**
 - Money does not write software. RSEs do. **We need to create attractive career plans.**
 - We need to make research software development attractive to students. **Academic recognition. Industry career paths.**
- **Anticipating the future in hardware development accelerates the porting process.**
 - **Blueprints** and **early access systems** both useful.
 - **Interaction with industry** is mutually beneficial.
- **Management, tools, and strategic initiatives, interaction and collegial behavior are important.**
 - Jira/Notion/[...] milestones and deliverables give projects and collaborative interactions a structure and timeline.
 - **Strategic focus groups, conferences,** and **meetings** bring experts together and **create collaboration.**
 - **Listen to the application needs. Value input and acknowledge collaborators.**

